

Sol'n HW4

9. Give a precise definition of what it means for the variable x to occur free as the i th symbol in the wff α . (If x is the i th symbol of α but does not occur free there, then it is said to occur *bound* there.)

(Sorry for this annoying problem!)

Suppose we've already defined $\text{len}(\varphi) = \text{length of } \varphi, \varphi \text{ a term or WFF.}$

Now we define $\Theta(i, \varphi) = \begin{cases} T & \text{if } x \text{ is free as } i^{\text{th}} \text{ symbol in } \varphi \\ F & \text{otherwise} \end{cases}$

Where φ is a WFF or term and $i \in \mathbb{N}$:

$$\Theta(i, c) = F, \quad c \text{ a constant symbol}$$

$$\Theta(i, v) = \begin{cases} T, & v = x \\ F, & v \neq x \end{cases}, \quad v \text{ a variable symbol}$$

$$\Theta(i, f(\tau_1, \dots, \tau_n)) = T \text{ provided } i = 2 + (j-1) + \sum_{m=1}^{j-1} \text{len}(\tau_m) + k, \text{ where}$$

$1 \leq j \leq n$ and $\Theta(k, \tau_j) = T$ for " f " \uparrow # of commas to j^{th} term

$$\Theta(i, (\tau_1 = \tau_2)) = T \text{ provided } \Theta(i-1, \tau_1) \text{ or } \Theta(i-1-\text{len}(\tau_1)-1, \tau_2) = T$$

$$\Theta(i, P(\tau_1, \dots, \tau_n)) = T \text{ provided } i = 2 + (j-1) + \sum_{m=1}^{j-1} \text{len}(\tau_m) + k, \text{ where}$$

$1 \leq j \leq n$ and $\Theta(k, \tau_j) = T$ for " P " \uparrow # of commas to j^{th} term

$$\Theta(i, (\neg \varphi)) = T \text{ provided } \Theta(i-2, \varphi) = T$$

$$\Theta(i, (\varphi \rightarrow \psi)) = T \text{ provided } \Theta(i-1, \varphi) = T \text{ or } \Theta(i-1-\text{len}(\varphi)-1, \psi) = T$$

$$\Theta(i, (\forall x \varphi)) = F \text{ if } v = x, = T \text{ if } v \neq x \text{ and } \Theta(i-2, \varphi)$$

OK, how to define $\text{len}(\varphi) = \text{length of } \varphi, \varphi \text{ a term or WFF:}$

$$\text{len}(c) = \text{len}(v) = 1, \quad c \text{ a constant symbol, } v \text{ a variable symbol}$$

$$\text{len}(f(\tau_1, \dots, \tau_n)) = 2 + n + \sum_{i=1}^n \text{len}(\tau_i)$$

$$\text{len}(\tau_1 = \tau_2) = 3 + \text{len}(\tau_1) + \text{len}(\tau_2)$$

$$\text{len}(P(\tau_1, \dots, \tau_n)) = 2 + n + \sum_{i=1}^n \text{len}(\tau_i)$$

$$\text{len}(\neg \varphi) = 3 + \text{len}(\varphi)$$

$$\text{len}(\varphi \rightarrow \psi) = 3 + \text{len}(\varphi) + \text{len}(\psi)$$

$$\text{len}(\forall x \varphi) = 2 + \text{len}(\varphi)$$

... and that's it! Whew!

2. Let $\mathcal{L} = \{P\}$ be the language with precisely one unary predicate symbol P . Suppose $A = \{1, 2, 3, \dots, N\}$ for some natural number N . How many different \mathcal{L} -structures are there with domain A ?

Every distinct $P_{\mathcal{A}} \subseteq A$ gives a different \mathcal{L} -structure. Thus the # of different \mathcal{L} -structures with universe A is $\text{Card}(\mathcal{P}(A)) = 2^n$

4. Let $\mathcal{L} = \{E\}$ be the language with precisely one binary predicate symbol. Consider the \mathcal{L} -structures $\mathfrak{A} = \{\mathbb{N}, <_{\mathbb{N}}\}$, $\mathfrak{B} = \{\mathbb{Q}, <_{\mathbb{Q}}\}$, and $\mathfrak{C} = \{\mathbb{Z}, <_{\mathbb{Z}}\}$. For each pair of these models, find a sentence true in one but not the other.

$\exists x \forall y E(x, y)$ is true in \mathfrak{A} , not in $\mathfrak{B}, \mathfrak{C}$.

$\forall x \forall y (E(x, y) \rightarrow (\exists z E(x, z) \wedge E(z, y)))$ is true in \mathfrak{B} , not in \mathfrak{C}

6. A model \mathfrak{A} for the language \mathcal{L} (from problem 5) is defined by

$$\mathfrak{A} = \langle \mathbb{N}, 0_{\mathbb{N}}, f_{\mathfrak{A}}, <_{\mathbb{N}} \rangle,$$

where $f_{\mathfrak{A}}(x) = x + 1$ for all $x \in \mathbb{N}$. For each of the formulas in 5(a)-(e) above, if the formula is a sentence state whether or not it holds in \mathfrak{A} .

The sentences are (a), (d), (e):

(a) $\forall x \exists y P(y, x)$ In \mathfrak{A} this asserts that every x is bigger than some y , but 0 is not bigger than any y , so false in \mathfrak{A}

(d) $\exists x (x = f(x))$ In \mathfrak{A} this says some x is fixed by f , but x never $= x + 1$ in \mathbb{N} , so false in \mathfrak{A}

(e) $\forall x P(x, f(x))$ Says for all x , $x < x + 1$, which is true in \mathfrak{A}

7. A model \mathfrak{B} for the language \mathcal{L} (from problem 5) is defined by

$$\mathfrak{B} = \langle \mathbb{Q}, 0_{\mathbb{Q}}, f_{\mathfrak{B}}, <_{\mathbb{Q}} \rangle,$$

where $f_{\mathfrak{B}}(x) = x/2$ for all $x \in \mathbb{Q}$. For each of the formulas in 5(a)-(e) above, if the formula is a sentence state whether or not it holds in \mathfrak{B} .

(a) $\forall x \exists y P(y, x)$ Every x in \mathbb{Q} is bigger than some y , so true in \mathfrak{B}

(d) $\exists x (x = f(x))$ Since $0 = 0/2$, this is true in \mathfrak{B}

(e) $\forall x P(x, f(x))$ If $x > 0$, $x \not< x/2$, so false in \mathfrak{B} .