When can $eX=1 \mod (p-1)(q-1)$ be solved? Itow? More generally, to solve $eX=1 \mod m$

Start with an assumption:

ASSUME 1= GCF(e,m) (That is, e,m relatively prime)
Use the Eudidean Algorithm to find integers r,s such That

v.e + S.m = 1

Then of corre m | s.m, so re=1 mod m

So we're done! (Just use this value of r.)

In The case of RSA, that's all we need; This r becomes

The private exponent of.

We can do This more generally: Suppose we want to solve The equation

ex=b mod m

ASSUME FIRST MAT GCF(e,m) | b. That is, if

g=GCF(e,m) Then b=gK for some integer K.

Again wing the Evolidean algorithm, find ris

such that

r.e + s·m = g

Then (multiply by K) (riche + (sichem = gik divisible = b. -

.. e. (rx) = b mod m, So a solution is rx.

Example Solve 90x = 18 mod 12

Use Euclidean Algorithm to find CCF(90,D): 90 = 12.7 + 6 12 = 6.2 + 0So 6 = GCF(90,12). Note 6/18, infact $\frac{18}{6} = 3$: This K (laker).

Now, 6 = 90 - 12.7 = 1.90 + (-1).12,

and 18 = 3.90 + (-1).12 (I just multiplied throughly 16 = 3)

for 90.3 = 18 mod 12So 90.3 = 18 mod 12

Finally a comment about the assumptions. In the exil example we assumed exim were relatively prime; in the exib modern example we assumed GCF(e,m) | b. What IF This isn't tree.

Suppose we have any solin of the equation ex=b mod m.

Then ex-b is a multiple of m, say (ex-b) = t m.

Newrik: b = tm-ex.

The CCF of m and e most divide both m and e, so divides the right-hand side (tm-ex), so divides b. That tells is that our assumption was necessary, that we can only solve ex=b mod m if CCF(e,m) | b. Bye for now.